JC08 Rec'd PCT/PTO 2 1 FEB 2001

FORM PTO-1390 (REV. 11-2000) U S DEPARTMENT OF COMMERCE PATENT AND TRADEMARK OFFICE

ATTORNEY 'S DOCKET NUMBER LKAR01US

TRANSMITTAL LETTER TO THE UNITED STATES DESIGNATED/ELECTED OFFICE (DO/EO/US) CONCERNING A FILING UNDER 35 U.S.C. 371

U S APPLICATION NO (If known, see 37 CFR 1 5 09/763289

INTERNATIONAL APPLICATION NO.	INTERNATIONAL FILING DATE	PRIORITY DATE CLAIMED
PCT/CN00/00028	17 FEBRUARY 2000	17 FEBRUARY 2000

TITLE OF INVENTION A SPREAD-SPECTRUM MULTIPLE ACCESS CODING METHOD WITH ZERO CORRELATION WINDOW

APPLICANT(S) FOR DO/EO/US DAOBEN LI

Applicant herewith submits to the United States Designated/Elected Office (DO/EO/US) the following items and other information:

1. ☒ This is a **FIRST** submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a **SECOND** or **SUBSEQUENT** submission of items concerning a filing under 35 U.S.C. 371.
3. ☐ This is an express request to begin national examination procedures (35 U.S.C. 371(f)). The submission must include items (5), (6), (9) and (21) indicated below.
4. ☐ The US has been elected by the expiration of 19 months from the priority date (Article 31).
5. ☒ A copy of the International Application as filed (35 U.S.C. 371(c)(2))
   a. ☒ is attached hereto (required only if not communicated by the International Bureau).
   b. ☐ has been communicated by the International Bureau.
   c. ☐ is not required, as the application was filed in the United States Receiving Office (RO/US).
6. ☒ An English language translation of the International Application as filed (35 U.S.C. 371(c)(2)).
   a. ☒ is attached hereto.
   b. ☐ has been previously submitted under 35 U.S.C. 154(d)(4).
7. ☐ Amendments to the claims of the International Aplication under PCT Article 19 (35 U.S.C. 371(c)(3))
   a. ☐ are attached hereto (required only if not communicated by the International Bureau).
   b. ☐ have been communicated by the International Bureau.
   c. ☐ have not been made; however, the time limit for making such amendments has NOT expired.
   d. ☐ have not been made and will not be made.
8. ☐ An English language translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 371 (c)(3)).
9. ☒ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)).
10. ☐ An English lanugage translation of the annexes of the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Items 11 to 20 below concern document(s) or information included:

11. ☐ An Information Disclosure Statement under 37 CFR 1.97 and 1.98.
12. ☐ An assignment document for recording. A separate cover sheet in compliance with 37 CFR 3.28 and 3.31 is included.
13. ☐ A FIRST preliminary amendment.
14. ☐ A SECOND or SUBSEQUENT preliminary amendment.
15. ☐ A substitute specification.
16. ☐ A change of power of attorney and/or address letter.
17. ☐ A computer-readable form of the sequence listing in accordance with PCT Rule 13ter.2 and 35 U.S.C. 1.821 - 1.825.
18. ☐ A second copy of the published international application under 35 U.S.C. 154(d)(4).
19. ☐ A second copy of the English language translation of the international application under 35 U.S.C. 154(d)(4).
20. ☐ Other items or information:

JC02 Rec'd PCT/PTO 21 FEB 2001

U.S. APPLICATION NO (if known, see 37 CFR 1.5) 09/763289	INTERNATIONAL APPLICATION NO PCT/CN00 00028	ATTORNEY'S DOCKET NUMBER LKAR01US

21. [X] The following fees are submitted:

BASIC NATIONAL FEE (37 CFR 1.492 (a) (1) - (5)):

	CALCULATIONS	PTO USE ONLY
Neither international preliminary examination fee (37 CFR 1.482) nor international search fee (37 CFR 1.445(a)(2)) paid to USPTO and International Search Report not prepared by the EPO or JPO **$1000.00**		
International preliminary examination fee (37 CFR 1.482) not paid to USPTO but International Search Report prepared by the EPO or JPO **$860.00**		
International preliminary examination fee (37 CFR 1.482) not paid to USPTO but international search fee (37 CFR 1.445(a)(2)) paid to USPTO **$710.00**		
International preliminary examination fee (37 CFR 1.482) paid to USPTO but all claims did not satisfy provisions of PCT Article 33(1)-(4) **$690.00**		
International preliminary examination fee (37 CFR 1.482) paid to USPTO and all claims satisfied provisions of PCT Article 33(1)-(4) **$100.00**		
ENTER APPROPRIATE BASIC FEE AMOUNT =	$ 1,000.00	
Surcharge of **$130.00** for furnishing the oath or declaration later than [ ] 20 [ ] 30 months from the earliest claimed priority date (37 CFR 1.492(e)).	$ 0.00	

CLAIMS	NUMBER FILED	NUMBER EXTRA	RATE		
				$	
Total claims	32 -20 =	12	x **$18.00**	$ 216.00	
Independent claims	0 -3 =	0	x **$80.00**	$ 0.00	
MULTIPLE DEPENDENT CLAIM(S) (if applicable)			+ **$270.00**	$	

TOTAL OF ABOVE CALCULATIONS =	$ 1,216.00	
[X] Applicant claims small entity status. See 37 CFR 1.27. The fees indicated above are reduced by 1/2. +	$ 608.00	
SUBTOTAL =	$ 608.00	
Processing fee of **$130.00** for furnishing the English translation later than [ ] 20 [ ] 30 months from the earliest claimed priority date (37 CFR 1.492(f)).	$ 0.00	
TOTAL NATIONAL FEE =	$ 608.00	
Fee for recording the enclosed assignment (37 CFR 1.21(h)). The assignment must be accompanied by an appropriate cover sheet (37 CFR 3.28, 3.31). **$40.00** per property +	$ 0.00	
TOTAL FEES ENCLOSED =	$ 608.00	
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a. [ ] A check in the amount of $ 608.00 to cover the above fees is enclosed.

b. [ ] Please charge my Deposit Account No. ____________ in the amount of $ __________ to cover the above fees. A duplicate copy of this sheet is enclosed.

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NOTE: Where an appropriate time limit under 37 CFR 1.494 or 1.495 has not been met, a petition to revive (37 CFR 1.137 (a) or (b)) must be filed and granted to restore the application to pending status.

SEND ALL CORRESPONDENCE TO

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[signature]
SIGNATURE

DAVID B. NEWMAN, JR.
NAME

30,966
REGISTRATION NUMBER

Attorney Docket: LKAR01US

DECLARATION CLAIMING SMALL ENTITY STATUS

I hereby declare that I qualify as an independent inventor/owner as defined in 37 C.F.R. § 1.9(c) for purposes of paying reduced fees under section 41(a) and (b) of Title 35, United States Code, to the Patent Office with regard to the invention entitled:

A SPREAD-SPECTRUM MULTIPLE ACCESS CODING METHOD
WITH ZERO CORRELATION WINDOW

and disclosed in a patent application for which an international patent application was filed 17 FEBRUARY 2000, having INTERNATIONAL APPLICATION NO. PCT/CN00/00028, and which is being filed concurrently herewith, by inventor

DAOBEN LI

I have not assigned, granted, conveyed or licensed and am under no obligation under contract or law to assign, grant, convey or license, any rights in the invention to any person who could not be classified as an independent inventor under 37 C.F.R. § 1.9(c) if that person had made the invention, or to any concern which would not qualify as a small business concern under 37 C.F.R. § 1.9(d) or a nonprofit organization under 37 C.F.R. § 1.9(e).

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I acknowledge the duty to file, in this application or patent, notification if any change in status resulting in loss of entitlement to small entity status prior to paying, or at the time of paying, the earliest of the issue fee or any maintenance fee due on the date on which status as a small entity is no longer appropriate. (37 C.F.R. § 1.28(b))

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under section 1001 of Title 18 of the United States Code, and that such willful false statements may jeopardize the validity of the application, any patent issuing thereon, or any patent to which this verified statement is directed.

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3 PRTS

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JC02 Rec'd PCT/PTO 21 FEB 2001

BACKGROUND OF THE INVENTION

This invention relates to a spread-spectrum and code-division-multiple-access (CDMA) wireless communication technology; in particular, the invention relates to a spread-spectrum multiple access coding method having high spectral efficiency with zero correlation window in a Personal Communication System (PCS).

DESCRIPTION OF THE RELEVANT ART

The growing popularity of personal communication services coupled with the scarcity of radio bandwidth resources has resulted in the ever-increasing demand for higher spectral efficiency in wireless communications. Spectral efficiency refers to the maximum number of subscribers that can be supported in a cell or sector under a given bandwidth allocation and transmission rate requirement. The unit of spectral efficiency is the total transmission rate per unit bandwidth within a given cell or sector. Obviously, the better the spectral efficiency is, the higher the system capacity will be.

Traditional wireless Multiple Access Control (MAC) schemes, such as Frequency Division Multiple Access (FDMA), Time Division Multiple Access (TDMA), result in system capacity that is limited by the time-bandwidth multiple. It is impossible to increase the number of supportable subscribers under these MAC schemes. For example, assume that the basic transmission rate

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of a subscriber is 1/T samples per second and the allocated bandwidth is B Hz. Then, the time-bandwidth multiple is BT, which is the maximum number of supportable subscribers. It is impossible to support more than BT subscribers in FDMA and TDMA systems.

The situation is completely different under Code Division Multiple Access (CDMA) scheme where the system capacity only depends on the Signal-to-Interference Ratio (SIR). Increasing the number of subscriber reduces the SIR, thus lowering the transmission rate. However, the subscriber will not be denied radio resource allocation. In other words, unlike FDMA and TDMA, CDMA does not have a hard upper bound (i.e. BT) on the number of supportable subscribers.

The capacity of a CDMA system depends on the interference level. As such, the ability to accurately control the interference level is critical to the performance and the successful operation of a CDMA system. There are four sources of interference in a CDMA system: the first type of interference (or noise) comes from various sources in the local environment, which cannot be control by the wireless communication system. The only way to alleviate this kind of interference is the use of low noise amplifier; the second type of interference is Inter-Symbol-Interference (ISI); the third type of interference is Multiple Access Interference (MAI) that is originated from

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other subscribers in the same cell; the forth type of interference is Adjacent Channel or Cell Interference (ACI) that is originated from other subscribers in the neighboring channel or cell. It is possible to reduce or eliminate ISI, MAI, and ACI by using higher performance codes.

In a CDMA system, each subscriber has his/her own unique identification code. In addition, the subscribers' spread-spectrum codes are orthogonal to each other. The orthogonality requirement is common to all multiple access schemes. If the communication channel is an ideal linear time and frequency non-dispersion system, and the system has high degree of synchronization, then the subscribers will stay orthogonal to each other. In reality, the communication channel is not ideal, and it is very difficult to achieve tight synchronization for communication channels with time and frequency dispersion. As a result, the ability to achieve orthogonality in a non-ideal communication channel with time and frequency dispersion is critical to the successful operation of CDMA systems.

It is commonly known that mobile communication channel is a typical random time varying channel, with random frequency dispersion (due to Doppler shift effect) and random time dispersion (due to multi-path transmission effect). The former results in the degradation in time selectivity of the received signal with unexpected fluctuation of the reception power level.

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The latter results in the degradation in frequency selectivity, which results in the unexpected variation in the reception level within each frequency component. This degradation results in reduced system performance and significantly lowers the system capacity. In particular, because of the time dispersion of the transmission channel (as a result of multi-path transmission), different signal paths do not arrive at the receiver at the same time. This results in the overlapping of neighboring symbols of the same subscriber and causes Inter Symbol Interference (ISI). On the other hand, the time dispersion of the channel worsens the multiple access interference. When the relative delay of signals of different subscribers are zero, any orthogonal code can achieve orthogonality. However, it is very hard to maintain orthogonality if the relative delay of signals of subscribers is not zero.

In order to reduce ISI, the auto-correlation of each subscriber's access codes must be an ideal impulse function that has all energy at the origin, nowhere else. To reduce the MAI, the cross-correlations between multiple access codes of different subscribers must be zero for any relative delay. In the terms of orthogonality, each access code must be orthogonal to itself with non-zero time delay. The access codes must be orthogonal to each other for any relative delay (including zero delay).

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For simplicity, the value of auto-correlation function at the origin is called the main lobe and the values of auto-correlations and cross-correlations at other points are called side lobes. The correlation functions of ideal multiple access codes should have zero side lobes everywhere. Unfortunately, it is proved in Welch theory that there does not exist any ideal multiple access codes in the field of finite elements and even in field of complex numbers. The claim that there do not exist ideal multiple access codes is called Welch bound. Especially, the side lobes of auto-correlation function and the side lobes of cross-correlation function are contradicted to each other; side lobes of one correlation function are small but the side lobes of the other correlation function become big. Furthermore, NASA had done brute force searching, by using computer, to search for all ideal codes. However, there had been no breakthrough. Since then, there has not been much research work on the search of the ideal multiple codes.

In fact, NASA had just searched for the good access codes in the Group codes and the Welch bound is true in the sub-fields of complex numbers. Beyond the field of complex numbers, the ideal codes could exist. For example, B. P. Schweitzer has found an approach to form ideal codes in his Ph.D thesis on "Generalized complementary code sets" in 1971. Later, Leppanen and Pentti (Nokia Telecommunication) extended Dr. Schweitser's results in the mixed TDMA and CDMA system. Their work has been

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granted a patent (No: 0600713A2; application number: 933095564). They broke the Welch bound in the high dimensional space. However, the utilization of frequency is very low and thus there is no practical value. There has not been any application of their invention in nearly 30 years. According to their invention, in a system of N multiple access codes, there requires at least N^2 basic codes. Each basic code has length at least N chips. That means it needs N^3 chips to support N addresses. For example, when N = 128, with 16QAM modulation, the coded spectral efficiency is only $\log_2 16 \times 128/128^3 = 2.441\times10^{-4}$ bits/Hz. The more access codes, the lower the utilization of the spectral efficiency. However, this coding methodology reminds us that ideal multiple access codes can be achieved via complementary code sets. However, we should avoid that the code length grows too fast with the required number of multiple access codes.

In addition, with technique of two-way synchronization, the relative time delay within each access code or between each other in a random time varying channel will not be greater than the maximum time dispersion of the channel plus the maximum timing error. Assuming that value is Δ second, so long as their correlation functions do not have any side lobes in a time interval $(-\Delta, \Delta)$, there are no MAI and ISI between the access codes. The time interval that possesses the above property is

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called "zero correlation window". It is obvious that the corresponding CDMA system will be ideal when the "zero correlation window" size is wider than the maximum time dispersion deviation of the channel (i.e. the time delays among multi-paths of the signal) plus the maximum timing error. At the same time, it is also true that the near-far effects are no longer effective. The well-known near-far effects is created by the overlapping of the side lobe of a signal source that is close to the base station receiver and the main lobe of a signal source that is far away from the base station receiver. The side lobe over-kills the main lobe, which causes high interference. The accurate, complicated and fast power control mechanism has to been used to overcome the near-far effects so that the energy of signals must be basically the same at the base station receiver. However, within the "zero correlation window" of the multiple access codes, there are no side lobes in the auto-correlation functions and cross-correlation functions under the working condition. The near-far effects no longer exist in the system. The complicated and fast power control mechanism will become less important and optional.

SUMMARY OF THE INVENTION

An object of the present invention is to provide a new coding method to create a series of spread-spectrum multiple access codes that have the "Zero Correlation Window" in their auto-correlation functions and cross-correlation functions. Due

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to the creation of the "zero correlation window", the fatal near-far effects in traditional CDMA radio communications is solved. The Multiple Access Interference (MAI) and the Inter-Symbol Interference (ISI) is extinguished. A high RF capacity radio system could be thus created based on the invention.

The spread-spectrum multiple access codes with "zero correlation window" according to the present invention has the following two properties: The auto-correlation functions are zero except at the origin where all energy resides. That means the multiple access codes are ideal in the sense that the access codes are orthogonal to themselves with any relative nonzero time delay. There exists a "zero correlation window" at the origin hereabout where the cross-correlation functions of spread-spectrum multiple access codes are zero everywhere inside the window. This means that the access codes are mutually orthogonal whenever the relative time delays are no more than the window size.

To achieve the above objective, the coding method of spread-spectrum multiple access codes with "Zero Correlation Window" according to the present invention includes the following steps:

Selecting a pair of basically orthogonal complementary code group (C1, S1), (C2, S2) with a code length N, in which the

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acylic auto-correlation and cross-correlation functions of code C and code S oppose each other but also complement each other except at the origin, after summarization of each other, the value of the auto-correlation and cross-correlation functions will be zero everywhere except at the origin;

Expanding the code length and code number of the pair of basically orthogonal complementary code group in a tree structure, according to the practically necessary maximum of subscriber access, the auto-correlation function of the expanded code group will be zero everywhere except at the origin, while the cross-correlation function will form a Zero Correlation Window around the origin with the size of the window ≥ 2N-1.

The width of Zero Correlation Window should be more than or equal to the maximum of relative time delay within each access code or between each other in the system. The maximum of relative time delay will be determined by the maximum time dispersion of the channel plus the maximum timing error.

When applying the formed spread-spectrum access codes in practical project, it should be ensured that code C only operate with code C (including itself and other codes), and code S only with code S (including itself and other codes). Therefore, using two orthogonal propagation channels that are synchronous fading, the above code C and code S can be transmitted

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respectively, and the same information bits can be loaded on modulation, and then summarize their output after despreading and demodulating. For the two orthogonal propagation channels, code C and code S can be modulated respectively on polarized waves orthogonal with each other, or code C and code S can be put in two time slots that will not overlap with each other after transmission.

The step of expanding the code length and code number of the pair of basically orthogonal complementary code group in a tree structure, according to the present invention, refers to:

If **(C1, S1)**, (**C2, S2**) were a pair of basically orthogonal complementary code group with code length N, then the two pairs of orthogonal complementary code group with each code length 2N can be generated in the following way:

```
                    C1  C2  S1  S2
               ┌─   C1  -C2  S1  -S2
C1  S1         │
C2  S2    ─────┤
               │
               └─   C2  C1  S2  S1
                    C2  -C1  S2  -S1
```

Wherein the values of auto-correlation functions of the orthogonal complementary code group formed on upper and lower trees after spread will be zero everywhere except at the origin, while the cross-correlation function will form a Zero

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Correlation Window around the origin with the size of the window $\geq 2N-1$.

The above spread can be kept going on in accordance with the tree structure so as to generate 2^{n+1} orthogonal complementary code groups with the code length $N2^n$ and the width of the zero correlation window $\geq 2N-1$, in which $n = 0,1,2,\ldots$ is the number of spread times.

The equivalent transformation can be made to the generated orthogonal complementary code group.

The pair of basically orthogonal complementary code group (C1, S1), (C2, S2), according to the present invention, refers to that the auto-correlation function and cross-correlation function is respectively the summation of acylic auto-correlation with cross-correlation functions between codes C, and the summation of acylic auto-correlation with cross-correlation functions between codes S.

The code length and the width of the zero correlation window of the pair of basically orthogonal
complementary code group can be spread in the following way:

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```
                    C1  C2  S1  S2
                 ┌─ C1  -C2  S1  -S2
C1  S1        ─┤
C2  S2          │
                 └─ C2  C1  S2  S1
                    C2  -C1  S2  -S1
```

Wherein if each code length of the pair of basically orthogonal complementary code group **(C1, S1)**, (**C2, S2**) is N, and the width of the zero correlation window is L, then each code length of the spread pair of basically orthogonal complementary code group will be 2N, while the width of the zero correlation window will be 2L+1.

When N = 2, the pair of basically orthogonal complementary ode group will be:

(++ , +-)

(-+ , --)

Wherein "+" means +1 and "-" means -1, while the width of the zero correlation window will be 3.

The above spread can be kept going on in accordance with the tree structure so as to generate 2^n pairs of orthogonal complementary code groups with the code length $N2^n$ and the width

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of the zero correlation window as $2^nL+2^{n-1}+2^{n-2}+2^{n-3}+\ldots+2^1+1$, in which n = 0,1,2, … is the number of spread times.

The equivalent transformation can be made to the generated basically orthogonal complementary code group.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a first schematic diagram of a generation tree of an orthogonal complementary code group with zero correlation window in the present invention;

FIG. 2 is a second schematic diagram of the generation tree of the orthogonal complementary code group with zero correlation window in the present invention; and

FIG. 3 is a schematic diagram of the generation tree of the basically orthogonal complementary code group in the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

The present invention is described with reference to the preferred embodiments and the drawings.

The coding steps of the present invention will be described hereinafter beginning with the basic code group with its code length 2 and the access number 2.

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Given two sets of codes of length 2, C Set: C1=(+, +), C2=(-, +) and S Set: S1=(+, -), S2=(-, -); wherein "+" means +1 and "-" means -1.

It is true that without any shift between each other (relative time delay), each pair of {C1, C2}, {S1, S2}, {C1, S1}, {C2, S2} are mutually orthogonal, i.e. their cross-correlation functions have zero value at the origin. However, with shift between each other (relative time delay), the mutual orthogonal property may not exist, i.e. the cross-correlation functions have non-zero values except the origin. See the following correlation tables for details. Table 1 shows the auto- and cross-correlation functions values of codes **C**1 and **C**2 with different shifts and Table 2 shows the auto- and cross-correlation values of codes **S**1 and **S**2 with different shifts.

Table 1 Correlation of the C Codes: C1=(+ +); C2 =(- +)

Time shift τ / Correlation	-1	0	1
$R_{c_1}(\tau)$	1	2	1
$R_{c_2}(\tau)$	-1	2	-1
$R_{c_1c_2}(\tau)$	1	0	-1

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Table 2 Correlation of the S codes: S1=(+ -); S2=(- -)

Time shift τ / Correlation	-1	0	1
$R_{s_1}(\tau)$	-1	2	-1
$R_{s_2}(\tau)$	1	2	1
$R_{s_1 s_2}(\tau)$	-1	0	1

It can be seen that both codes are not ideal. However, when adding these two tables together, the codes become ideal (See Table 3).

Now Define auto-correlation functions

$$R_1(\tau) \stackrel{\Delta}{=} R_{c_1}(\tau)+R_{c_2}(\tau),\ R_2(\tau) \stackrel{\Delta}{=} R_{c_2}(\tau)+R_{s_2}(\tau),$$

and cross-correlation functions

$$R_{12}(\tau) \stackrel{\Delta}{=} R_{c_1 c_2}(\tau)+R_{s_1 s_2}(\tau).$$

With the above new definition of correlation functions, i.e. the new correlation functions (including auto-and cross-correlation functions) are summation of the correlation functions of C codes and the correlation functions of S codes,

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the values of auto- and cross-correlation functions of the codes one and codes two become ideal.

Such codes C and S can be called "complementary orthogonal" if C and S are ideal under the new definition of correlation functions $R_1(\tau)$, $R_2(\tau)$, and $R_{12}(\tau)$, i.e. their correlation functions are opposed and complemented to each other except the origin. The above **C** and **S** code sets can be, for convenience, expressed as (C1, S1) = (+ +, + -) and (C2, S2) = (- +, - -).

Table 3 shows the correlation functions of the complementary orthogonal codes.

Table 3 Correlation of C and S codes (C1, S1) = (+ +; + -); (C2, C2)= (- +; - -)

Time shift τ / Correlation	-1	0	1
$R_1(\tau) \triangleq R_{c_1}(\tau) + R_{s_1}(\tau)$	0	4	0
$R_2(\tau) \triangleq R_{c_2}(\tau) + R_{s_2}(\tau)$	0	4	0
$R_{12}(\tau) \triangleq R_{c_1c_2}(\tau) + R_{s_1s_2}(\tau)$	0	0	0

There is only one basic form for the orthogonal complementary code group with the number of access code 2 and each code length 2. It is proven that the C set of codes C1 = (+ +), C2 = (- +) and the S set of codes: S1 = (+ -), S2 = (- -) are the basic form of complementary orthogonal codes of length 2. Other forms can be derived from re-ordering of C1 and C2, S1 and S2, swapping C and S, rotation, order reverse, interleaving polarity, and alternative negation etc without any substantial differences. It should be noted that only the operation of code C with code C and code S with code S should take place when making the operation of correlation or matching filtering. Code C and code S will not encounter on operation.

For longer code, for example, the orthogonal complementary code group with the number of access code 2 and each code length 4 can be derived from the above basically orthogonal complementary code group.

One of the generation methods is:

Let

(C1', S1') = (C1 C2, S1 S2);

(C2', S2') = (**C**1 - **C**2, S1 - S2);

Wherein C1' means the concatenation of original code **C**1 and **C**2;

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C2′ means the concatenation of **C**1 and the negation of the **C**2. Same operations could be applied to S1′ and S2′.

They can be expressed as:

$$(C1', S1') = (+ + - +, + - - -);$$
$$(C2', S2') = (+ + + -, + - + +);$$

Table 4 shows the orthogonal complementary correlation functions of the new code group. It can be seen that the complementary auto-correlation function and cross-correlation function are all ideal.

The other way is reversing the order of the codes, that is:

$$(C1'', S1'') = (C2C1, S2S1) = (- + + +, - - + -)$$
$$(C2'', S2'') = (C2 - C1, S2 - S1) = (- + - -, - - - +)$$

The complementary auto-correlation function and cross-correlation function are also ideal. The orthogonal complementary correlation functions of the new code group are the same with those of the above code group. (See Table 4)

Table 4: The Orthogonal Complementary Correlation Functions (each code length is $2^2 = 4$):

$$(C1', S1') = (+ + - +, + - - -);$$
$$(C2', S2') = (+ + + -, + - + +);$$

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or

$(C1'', S1'') = (- + + +, - - + -)$

$(C2'', S2'') = (- + - -, - - - +)$

Time shift τ / Correlation	-3	-2	-1	0	1	2	3
$R_1(\tau)=R_{c_1}(\tau)+R_{s_1}(\tau)$	0	0	0	8	0	0	0
$R_2(\tau)=R_{c_2}(\tau)+R_{s_2}(\tau)$	0	0	0	8	0	0	0
$R_{12}(\tau)=R_{c_1c_2}(\tau)+R_{s_1s_2}(\tau)$	0	0	0	0	0	0	0

With this way going on, the orthogonal complementary code group with the number of access code 2 and each code length 2^n ($n = 1, 2\ldots$) can be obtained. It can be proved that their auto-correlation and cross-correlation functions are all ideal. Although the auto-correlation and cross-correlation functions of the access codes formed by this coding method, however, are ideal, the number of the access codes is only 2. It is apparent that two access codes are too small for a CDMA communications

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system. In practice, it is required that the number of the orthogonal access codes be as many as possible under the condition of given code length, while their auto-correlation and cross-correlation functions are not necessarily ideal everywhere. It is desirable that there is a zero correlation window around the origin that can meet the needs.

In fact, renumbering and arranging the above four complementary code groups with each code length 4, the result can be as follows:

(C1, S1) = (+ + − +, + − − −); (C2, S2) = (+ + + −, + − + +)

(C3, S3) = (− + + +, − − + −); (C4, S4) = (− + − −, − − − +)

Table 5 shows the correlation functions of the complementary code group.

Table 5: The Correlation Matrix of Codes (each code length is $2^2 = 4$):

(C1, S1) = (+ + − +, + − − −); (C2, S2) = (+ + + −, + − + +)

(C3, S3) = (− + + +, − − + −); (C4, S4) = (− + − −, − − − +)

Time shift τ / Correlation	-3	-2	-1	0	1	2	3
$R_1(\tau) \stackrel{\Delta}{=} R_{c_1}(\tau) + R_{s_1}(\tau)$	0	0	0	8	0	0	0
$R_2(\tau) \stackrel{\Delta}{=} R_{c_2}(\tau) + R_{s_2}(\tau)$	0	0	0	8	0	0	0
$R_3(\tau) \stackrel{\Delta}{=} R_{c_3}(\tau) + R_{s_3}(\tau)$	0	0	0	8	0	0	0
$R_4(\tau) \stackrel{\Delta}{=} R_{c_4}(\tau) + R_{s_4}(\tau)$	0	0	0	8	0	0	0
$R_{12}(\tau) \stackrel{\Delta}{=} R_{c_1c_2}(\tau) + R_{s_1s_2}(\tau)$	0	0	0	0	0	0	0
$R_{34}(\tau) \stackrel{\Delta}{=} R_{c_3c_4}(\tau) + R_{s_3s_4}(\tau)$	0	0	0	0	0	0	0
$R_{13}(\tau) \stackrel{\Delta}{=} R_{c_1c_3}(\tau) + R_{s_1s_3}(\tau)$	0	4	0	0	0	4	0
$R_{14}(\tau) \stackrel{\Delta}{=} R_{c_1c_4}(\tau) + R_{s_1s_4}(\tau)$	0	-4	0	0	0	4	0
$R_{23}(\tau) \stackrel{\Delta}{=} R_{c_2c_3}(\tau) + R_{s_2s_3}(\tau)$	0	4	0	0	0	-4	0
$R_{24}(\tau) \stackrel{\Delta}{=} R_{c_2c_4}(\tau) + R_{s_2s_4}(\tau)$	0	-4	0	0	0	-4	0

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Wherein (C1, S1) and (C2, S2), (C3, S3) and(C4, S4) are the pair of orthogonal complementary code group with ideal property respectively, but the cross-correlation functions between groups are not ideal. For example, $R_{13}(\tau)$ and $R_{14}(\tau)$, $R_{23}(\tau)$ and $R_{24}(\tau)$ are not zero everywhere, but there is a zero correlation window with the size of 3 chips wide. Thus, an orthogonal complementary code group with the number of access codes 4, each code length 4, and a zero correlation window can be obtained. The reason that the size of the zero correlation window is 3 is because the above four orthogonal complementary code groups are all composed of the basically orthogonal complementary code group with each code length 2, i.e. (C1, S1) = (+ +, + -) and (C2, S2) = (- +, - -), while the basic code group has only three status of time shift, i.e. -1, 0, and 1, because of each code length 2. In the ideal cases can only zero correlation window with the size of 3 be obtained.

To generate a wide window of zero correlation, the **C**1 and **S**1 codes are required to increase their sizes. For example, the code length can be 4. There are two pairs of completely orthogonal basic complementary code group with each code length 4.

They are: (+ + - +, + - - -), (+ + + -, + - + +), and (- + ++,

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- - + -), (-   +--, - - - +).

Supposing that the first pair of code group is the original orthogonal complementary code group, four pairs of orthogonal complementary code group with each code length 8 can be generated following the aforementioned methods.

They are: (C1, S1) = (+ + - + + + + -, + - - - + - + +); (C2, S2) = (+ + - + - - - +, + - - - - + - -); and (C3, S3) = (+ + + - + + - +, + - + + + - - -); (C4, S4) =(+ + + - - - + -, + - + + - + + +).

It can be expected that the size of their zero correlation window is 7 chips wide.

The correlation functions of these orthogonal complementary codes group are presented in the following matrix of Table 6:

Table 6 Correlation Matrix of codes (each code length $2^3 = 8$):

(C1, S1) = (+ + - + + + + -, + - - - + - + +);

(C2, S2) = (+ + - + - - - +, + - - - - + - -);

(C3, S3) = (+ + + - + + - +, + - + + + - - -);

(C4, S4) =(+ + + - - - + -, + - + + - + + +)

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Tine shift τ / Correlation	-7	-6	-5	-4	-3	-2	-1	0	1	2	3	4	5	6	7
$R_1(\tau) \overset{\Delta}{=} R_{c_1}(\tau) + R_{s_1}(\tau)$	0	0	0	0	0	0	0	16	0	0	0	0	0	0	0
$R_2(\tau) \overset{\Delta}{=} R_{c_2}(\tau) + R_{s_2}(\tau)$	0	0	0	0	0	0	0	16	0	0	0	0	0	0	0
$R_3(\tau) \overset{\Delta}{=} R_{c_3}(\tau) + R_{s_3}(\tau)$	0	0	0	0	0	0	0	16	0	0	0	0	0	0	0
$R_4(\tau) \overset{\Delta}{=} R_{c_4}(\tau) + R_{s_4}(\tau)$	0	0	0	0	0	0	0	16	0	0	0	0	0	0	0
$R_{12}(\tau) \overset{\Delta}{=} R_{c_1c_2}(\tau) + R_{s_1s_2}(\tau)$	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
$R_{34}(\tau) \overset{\Delta}{=} R_{c_3c_4}(\tau) + R_{s_3s_4}(\tau)$	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
$R_{13}(\tau) \overset{\Delta}{=} R_{c_1c_3}(\tau) + R_{s_1s_3}(\tau)$	0	0	0	8	0	0	0	0	0	0	0	8	0	0	0
$R_{14}(\tau) \overset{\Delta}{=} R_{c_1c_4}(\tau) + R_{s_1s_4}(\tau)$	0	0	0	-8	0	0	0	0	0	0	0	8	0	0	0
$R_{23}(\tau) \overset{\Delta}{=} R_{c_2c_3}(\tau) + R_{s_2s_3}(\tau)$	0	0	0	8	0	0	0	0	0	0	0	-8	0	0	0
$R_{24}(\tau) \overset{\Delta}{=} R_{c_2c_4}(\tau) + R_{s_2s_4}(\tau)$	0	0	0	-8	0	0	0	0	0	0	0	-8	0	0	0

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It is observed that two pairs of four new orthogonal complementary codes groups can be obtained from one pair of orthogonal complementary codes groups, with each code length doubled. Four pairs of eight orthogonal complementary codes groups can be further derived from these two pairs of four orthogonal complementary codes groups, and then, analogically in this way, eight pairs of sixteen orthogonal complementary codes groups can bé derived, …, wherein the auto-correlation functions of each codes group and the cross-correlation functions between inside codes groups are all ideal, while the cross-correlation functions of the codes groups between pairs have a zero correlation window with its size depending on the original orthogonal complementary code group. The process can be illustrated by some drawing of generation tree. Fig. 1 shows one of such generation tree, Fig. 2 is another one. There are many others of generation trees; the relations between them are equivalent transformation. Such transformation does not change the size of zero correlation windows. However, it sometimes changes the value of side lobes and their distribution outside the " zero correlation window".

FIG. 3 shows a basic pair of complementary code group which will be used in the actual coding process of multiple access codes. In Fig. 3, all pairs of code group in "<>" are basic pair of orthogonal complementary code group without any side

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lobes for their complementary auto-correlation functions and cross-correlation functions, that is to say, totally ideal. It should be noted that FIG. 3 shows only a pair of basically orthogonal complementary code group; there are still many ways of equivalent transformations, such as swapping the order of up and down or left and right, reversing the order of forwards and backwards, making alternately negation, rotating in complex plane, etc, in which equivalent pair of basically orthogonal complementary code group can be obtained with completely ideal auto-correlation and cross-correlation functions.

The construction process of the spread spectrum access codes according to the present invention will be described in detail below.

Firstly, determine the required size of zero correlation windows according as the propagation conditions of the applied system, the basic spread spectrum code bit rate (referred to as Chip Rate in terms of engineering, calculated as MCPS) used by the system, and the maximum timing error in the system.

Secondly, according to the required size of zero correlation window, select any pair of basically orthogonal complementary code group with its size of zero correlation

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window greater than or equal to the required window size as the original orthogonal complementary code group, and refer to it as (C1, S1), (C2, S2).

Then, determine the required maximum number of subscriber accesses according to the actual number of subscribers, and spread the selected original pair of basically orthogonal complementary code group as the origin of FIG. 2 or FIG. 3 in the tree view.

The number of spreading stages in FIG. 2 or FIG. 3 is dependent on the required maximum number of subscribers. For example, when the number of the required maximum number of subscribers is 120, because of $2^7 = 128 \geq 120$, then the required number of spreading stages is 7, while the $2^7 = 128$ group of codes in the 7^{th} stage of FIG. 2 or FIG. 3 can be the selected multiple access codes. At this time, the actual maximum number of subscriber accesses is 128, it is larger than 120, the required number of subscribers, and meets the needs completely.

In the practice of engineering, sometimes more mutations or variations of the access codes are needed. It needs to make equivalent transformation for the generated multiple access

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codes. The types of such transformations are so many that enumeration one by one is not necessarily. Here give the most common of equivalent transformations as follows:

Swapping the position of code C and code S.

Swapping the positions of C1 and C2 and S1 and S2 simultaneously.

Making negation to the order of codes.

Making negation to each code bit.

Interlacing the polarity of each code bit: for example, for (+ + - +, + - - -), (+ + + -, + - + +), interlace the polarity of each code bit, that is to say, the polarity of the odd code bits, such as the first, the third bit, etc, will remain unchanged, while the polarity of the even code bits, such as the second, the fourth bit etc, will change. So (+ - - -, + + - +), (+ - + +, + + + -) will result from this transformation. In like manner, the polarity of the odd code bits can be changed, while the polarity of the even code bits unchanged.

Rotating each code bit in complex plane: for example, by rotating in turn each code bit of (+ + - +, + - - -), (+ + + -, + - + +) at α angular degree, the following result will be obtained:

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$$(e^{j\varphi_{c_1}} e^{j(\varphi_{c_1}+\alpha)} - e^{j(\varphi_{c_1}+2\alpha)} e^{j(\varphi_{c_1}+3\alpha)}, e^{j\varphi_{s_1}} - e^{j(\varphi_{s_1}+\alpha)} - e^{j(\varphi_{s_1}+2\alpha)} - e^{j(\varphi_{s_1}+3\alpha)}$$

$$(e^{j\varphi_{c_2}} e^{j(\varphi_{c_2}+\alpha)} e^{j(\varphi_{c_2}+2\alpha)} - e^{j(\varphi_{c_2}+3\alpha)}, e^{j\varphi_{s_2}} - e^{j(\varphi_{s_2}+\alpha)} e^{j(\varphi_{s_2}+2\alpha)} e^{j(\varphi_{s_2}+3\alpha)}$$

Here $\varphi_{c_1}, \varphi_{c_2}, \varphi_{s_1}$ and φ_{s_2} can be any initial angular degree. It can be proven that the properties of auto-correlation and cross-correlation functions of each resultant access code are still unchanged after rotating transformation. However, the side lobes outside "zero correlation window" are relating to the rotating angular degree (being narrower or changing polarity). The aforementioned basically orthogonal complementary code group can be deemed as the code group with zero rotating angular degree.

Selecting properly the different rotating angular degree can make the rotated code groups orthogonal between them, i.e. multi groups of orthogonal codes can be generated from one group of orthogonal codes. This will be very convenient for the engineering application, especially when the code length is a little bit longer, sometimes the result will be so wonderful that it could meet various of actual needs of engineering, such as networking configuration, handoff/handovers, as well as the

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enhancement of RF capacity, etc.

Making transformation in the generation tree: for example, FIG. 3 is a kind of equivalent transform of FIG. 2, i.e. by moving all C1 codes and S1 codes to the left, C2 codes and S2 codes to the right in the corresponding C code and S code position; and interlacing, in certain rules, the code bits of C code and S code in the resulted multiple access codes groups, or changing the polarity arrangement, etc. In Mathematics, such transformation is called equivalent transformation. There are a lot of equivalent transforms that are impossible to enumerate one by one.

When applying the said formed spread spectrum access codes in practical project, it should be ensured that code C only operate with code C (including itself and other codes), and code S only with code S (including itself and other codes). Code C is never allowed to encounter code S. Therefore; the special parting measures should be taken in the actual application. For example, code C and code S can be modulated respectively on polarized waves (horizontal and vertical polarized waves, laevorotation and dextrorotation polarized waves) orthogonal with each other. Another example, code C and code S can be put in two time slots that will not overlap with each other after transmission. Because the propagation channels will change

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randomly with time, the channel properties within the two polarized waves and two time slots should be kept synchronous in the propagation process to ensure the complementarity. In terms of engineering, their fading should be synchronous. This means that when parting by polarization, the frequency channel without depolarization that can ensure the orthogonal polarized waves fading synchronously and corresponding measures should be used; when parting by time division, it should be ensured that the gap between two time slots is far less than the correlation time of channel; when using other parting methods, the synchronous fading should also be ensured.

Because code C and code S should be parted when propagation, and in the meantime, to utilize their complementarity, it is obvious that the data bits modulated on them should be identical, while the outputs after de-spreading and demodulation of code C and code S should be added together.

The coding method of the present invention presents a linear relation, because the total required number of code bits is only in direct proportion to the required number of accesses (about twofold). It moves forwards more creative step compared with the results of Dr. B.P. Schweitzer , Leppanen and Pentti. In their methods, the total required number of code bits is a

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cube relation with the required number of accesses. Therefore, it can be said that using the CDMA system according to the present invention will have much higher spectrum efficiency.

The present invention has been fully verified by computer simulation for four years. Under the same conditions, such as propagation fading, widening of multipath transmission, system bandwidth, subscriber transmission rate, and frame structure, as those of the first commercial CDMA standard in the world, i.e. IS-95, the spectrum efficiency of the system, when using the multiple access code system of the present invention, will be at least sixfold as that of IS-95.

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I CLAIM:

1. A construction method of the spreading spectrum multiple access codes with zero correlation window, wherein the method includes the following steps:

selecting a pair of basically orthogonal complementary code group (C1, S1), (C2, S2) with each code length as N chips, in which the acylic auto-correlation and cross-correlation functions of code C and code S oppose each other but also complement each other except at the origin, the values of auto-correlation and cross-correlation functions after summarization are zero except at the origin; and

based on the actually required maximum number of subscriber accesses, spreading the code length and code number of the basically orthogonal complementary code group in a tree structure, the values of auto-correlation functions of the spreaded code group are zero except at the origin, while the cross-correlation functions form a zero correlation window about the origin, with the window size greater than or equal to 2N-1.

2. The construction method of the spreading spectrum multiple access codes according to the claim 1, wherein the size of the zero correlation window is $\geq$ the maximum relative time delay inside each access code of the system or between them, the maximum relative time delay is dependent on the summation of the

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maximum time dispersion of the channel and the timing error of the system.

3. The construction method of the spreading spectrum multiple access codes according to the claim 1, wherein the above code C and code S are transmitted respectively by using two orthogonal and fading synchronously transmission channels, and carrying the same data bits when modulation, while the outputs are added together after de-spreading and demodulation.

4. The construction method of the spreading spectrum multiple access codes according to the claim 1, wherein the spreading the code length and code number of the basically orthogonal complementary code group in a tree structure refers to:

If (C1, S1), (C2, S2) is a pair of basically orthogonal complementary code group with code length N, the two pairs of orthogonal complementary code group with each code length 2N can be generated in the following way:

```
                 C1 C2 S1 S2
             ┌── C1 -C2 S1 -S2
C1 S1        │
C2 S2    ────┤
             │
             └── C2 C1 S2 S1
                 C2 -C1 S2 -S1
```

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wherein the values of auto-correlation functions of the orthogonal complementary code group formed on upper and lower trees after spread will be zero everywhere except at the origin, while the cross-correlation functions will form a Zero Correlation Window around the origin with the size of the window greater than or equal to 2N-1.

5. The construction method of the spreading spectrum multiple access codes according to the claim 4, wherein the above spread can be kept going on in accordance with the tree structure so as to generate 2^{n+1} orthogonal complementary code groups with the code length $N2^n$ and the width of the zero correlation window greater than or equal to 2N-1, in which n = 0,1,2, … is the number of spread times.

6. The construction method of the spreading spectrum multiple access codes according to the claim 4 or 5, wherein equivalent transformation can be applied to the resultant orthogonal complementary code group.

7. The construction method of the spreading spectrum multiple access codes according to the claim 6, wherein the equivalent transformation can be swap of the forward and backward position of the resultant code group.

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8. The construction method of the spreading spectrum multiple access codes according to the claim 6, wherein the equivalent transformation can be swap of the up and down position of the resultant code group.

9. The construction method of the spreading spectrum multiple access codes according to the claim 6, wherein the equivalent transformation can be negation of code order of each code.

10. The construction method of the spreading spectrum multiple access codes according to the claim 6, wherein the equivalent transformation can be interlacement of polarity of each code bit.

11. The construction method of the spreading spectrum multiple access codes according to the claim 6, wherein the equivalent transformation can be rotation of each code bit in complex plane in a sequence or without sequence.

12. The construction method of the spreading spectrum multiple access codes according to the claim 6, wherein the transformation can be any equivalent transformation that is proven in mathematics.

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13. The construction method of the spreading spectrum multiple access codes according to the claim 1, wherein the pair of basically orthogonal complementary code group (C1, S1), (C2, S2) refers to that the auto-correlation function and cross-correlation function is respectively the summation of acylic auto-correlation with cross-correlation functions between codes C, and the summation of acylic auto-correlation with cross-correlation functions between codes S.

14. The construction method of the spreading spectrum multiple access codes according to the claim 13, wherein the code length and the width of the zero correlation window of the pair of basically orthogonal complementary code group can be spread in the following way:

```
                    C1 C2 S1 S2
                ┌── C1 -C2 S1 -S2
C1 S1           │
C2 S2       ────┤
                │
                └── C2 C1 S2 S1
                    C2 -C1 S2 -S1
```

Wherein if each code length of the pair of basically orthogonal complementary code group (C1, S1), (C2, S2) is N, and the width of the zero correlation window is L, then each

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code length of the spread pair of basically orthogonal complementary code group will be 2N, while the width of the zero correlation window will be 2L+1.

15. The construction method of the spreading spectrum multiple access codes according to the claim 14, wherein when N = 2, the pair of basically orthogonal complementary code group will be:

(++ , +-)

(-+ , --)

Wherein "+" means +1 and "-" -1, while the width of the zero correlation window will be 3.

16. The construction method of the spreading spectrum multiple access codes according to the claim 14 or 15, wherein the above spread can be kept going on in accordance with the tree structure so as to generate 2^n pairs of orthogonal complementary code groups with the code length $N2^n$ and the width

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of the zero correlation window as $2^nL+2^{n-1}+2^{n-2}+2^{n-3}+\ldots+2^1+1$, in which n = 0,1,2, … is the number of spread times.

17. The construction method of the spreading spectrum multiple access codes according to the claim 16, wherein the equivalent transformation can be applied to the resultant basically orthogonal complementary code group.

18. The construction method of the spreading spectrum multiple access codes according to the claim 17, wherein the equivalent transformation can be swap of the forward and backward position of the resultant code group.

19. The construction method of the spreading spectrum multiple access codes according to the claim 17, wherein the equivalent transformation can be swap of the up and down position of the resultant code group.

20. The construction method of the spreading spectrum multiple access codes according to the claim 17, wherein the equivalent transformation can be negation of code order of each code.

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21. The construction method of the spreading spectrum multiple access codes according to the claim 17, wherein the equivalent transformation can be interlacement of polarity of each code bit.

22. The construction method of the spreading spectrum multiple access codes according to the claim 17, wherein the equivalent transformation can be rotation of each code bit in complex plane in a sequence or without sequence.

23. The construction method of the spreading spectrum multiple access codes according to the claim 17, wherein the transformation can be any equivalent transformation that is proven in Mathematics.

24. The construction method of the spreading spectrum multiple access codes according to the claim 3, wherein the orthogonal and fading synchronously transmission channel refers to the orthogonal polarized wave.

25. The construction method of the spreading spectrum multiple access codes according to the claim 3, wherein the orthogonal and fading synchronously transmission channel is the time slots without overlap to each other.

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26. The construction method of the spreading spectrum multiple access codes according to the claim 1, wherein one code or multiple access codes can be allocated based on the needs of the different data rate and services of each subscriber to actualize the different quality of priority level services.

27. The construction method of the spreading spectrum multiple access codes according to the claim 1, wherein the required spreading spectrum access codes can be adaptively generated based on the zero correlation window required by the different propagation modes, different number of subscribers, and the needs of different data rate as well as services, so that there are no inter-signal interference (ISI) and multi access interference (MAI) in the corresponding spreading spectrum CDMA system.

28. The construction method of the spreading spectrum multiple access codes according to the claim 1, wherein the resultant multiple access codes by the equivalent transformation can be used to meet the needs of network configuration, handoff and enhancement of system capacity, etc in cellular mobile or fixed point to multi points wireless telecommunications system.

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29. The construction method of the spreading spectrum multiple access codes according to the claim 1, wherein coding can be made, as one of the complex codes, by using complex codes.

30. The construction method of the spreading spectrum multiple access codes according to the claim 1, wherein it can be applied to any TD/CDMA, FD/CDMA, WD/CDMA, SD/CDMA or CDMA communications system.

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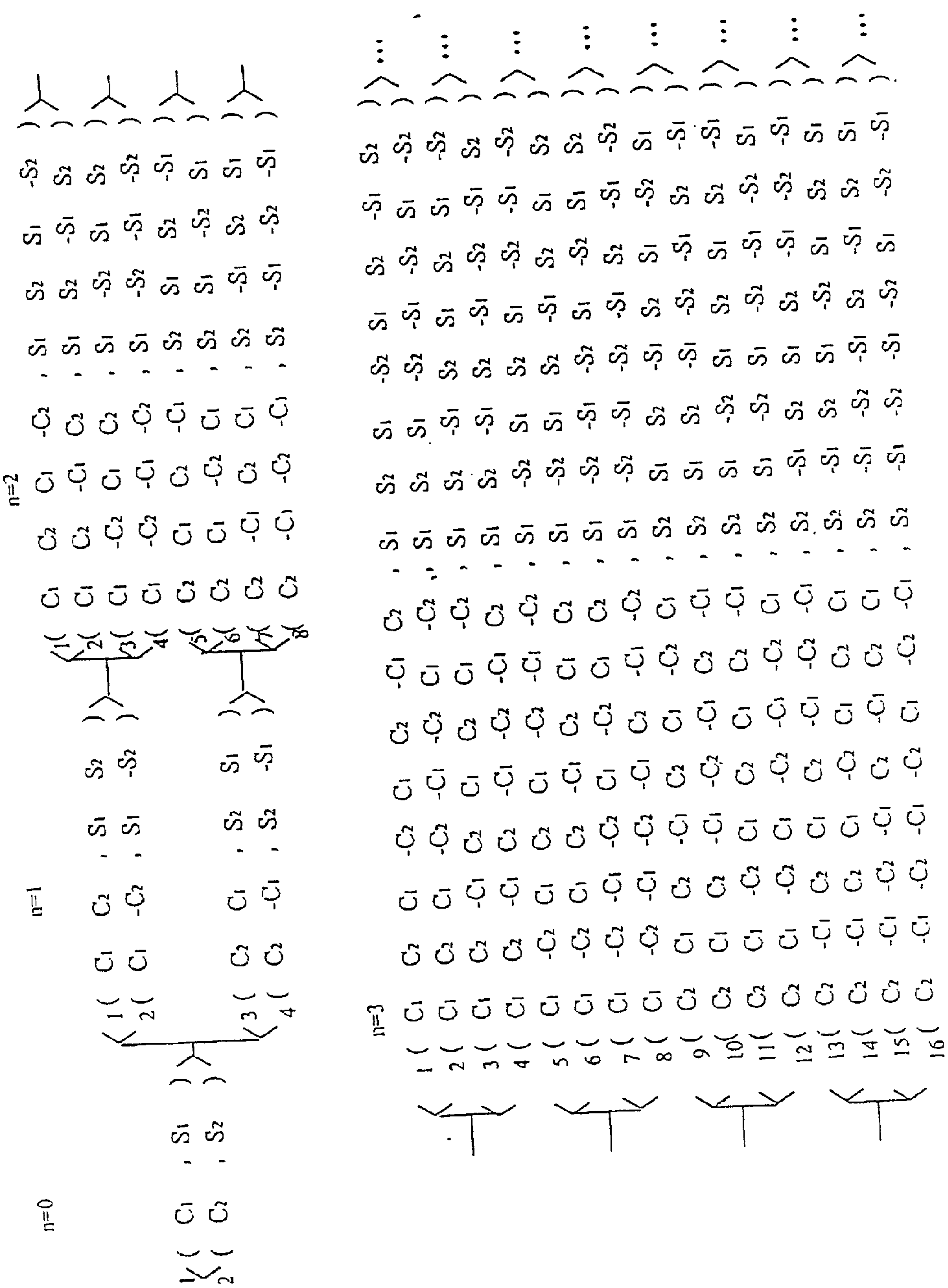

図1

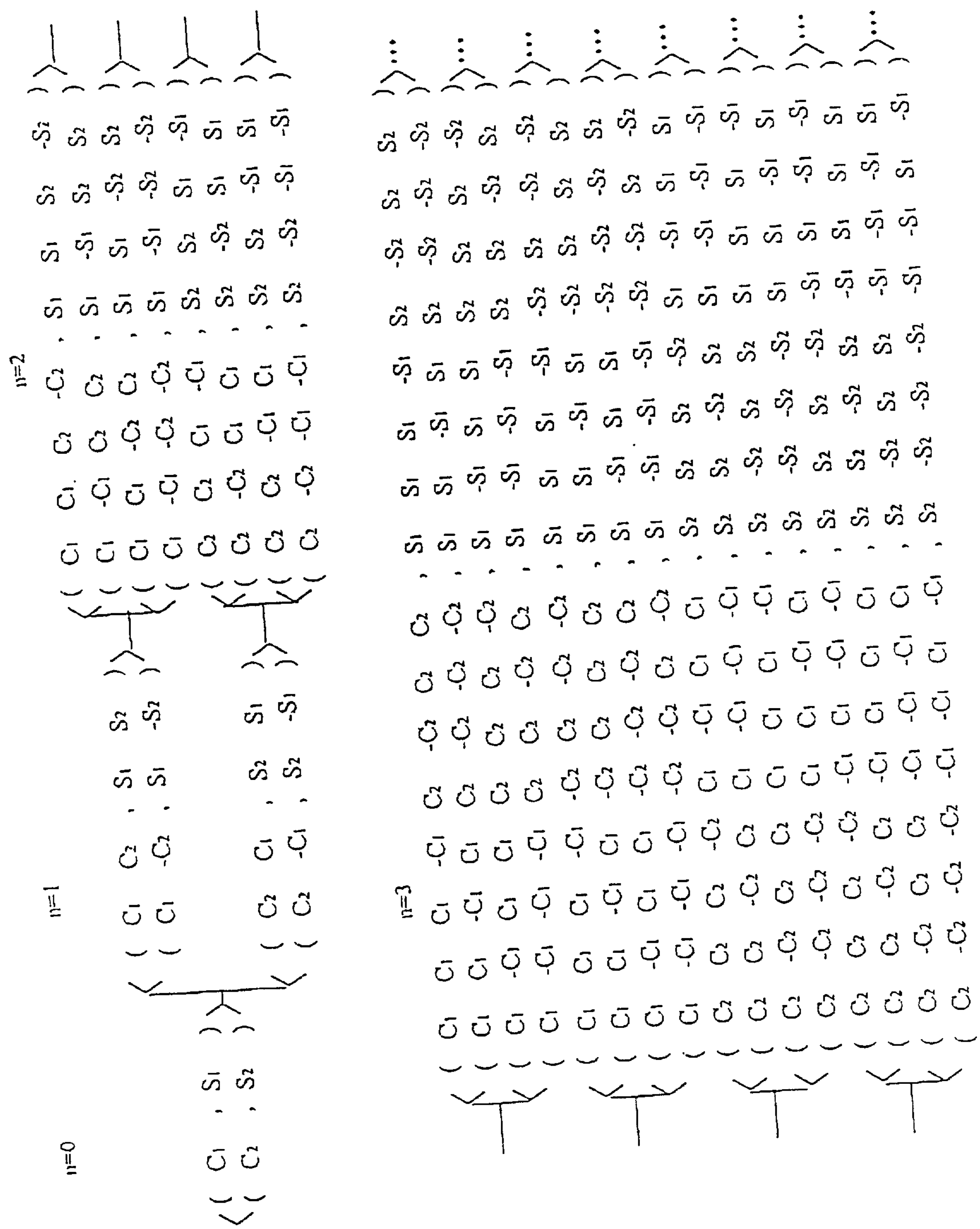

图 2

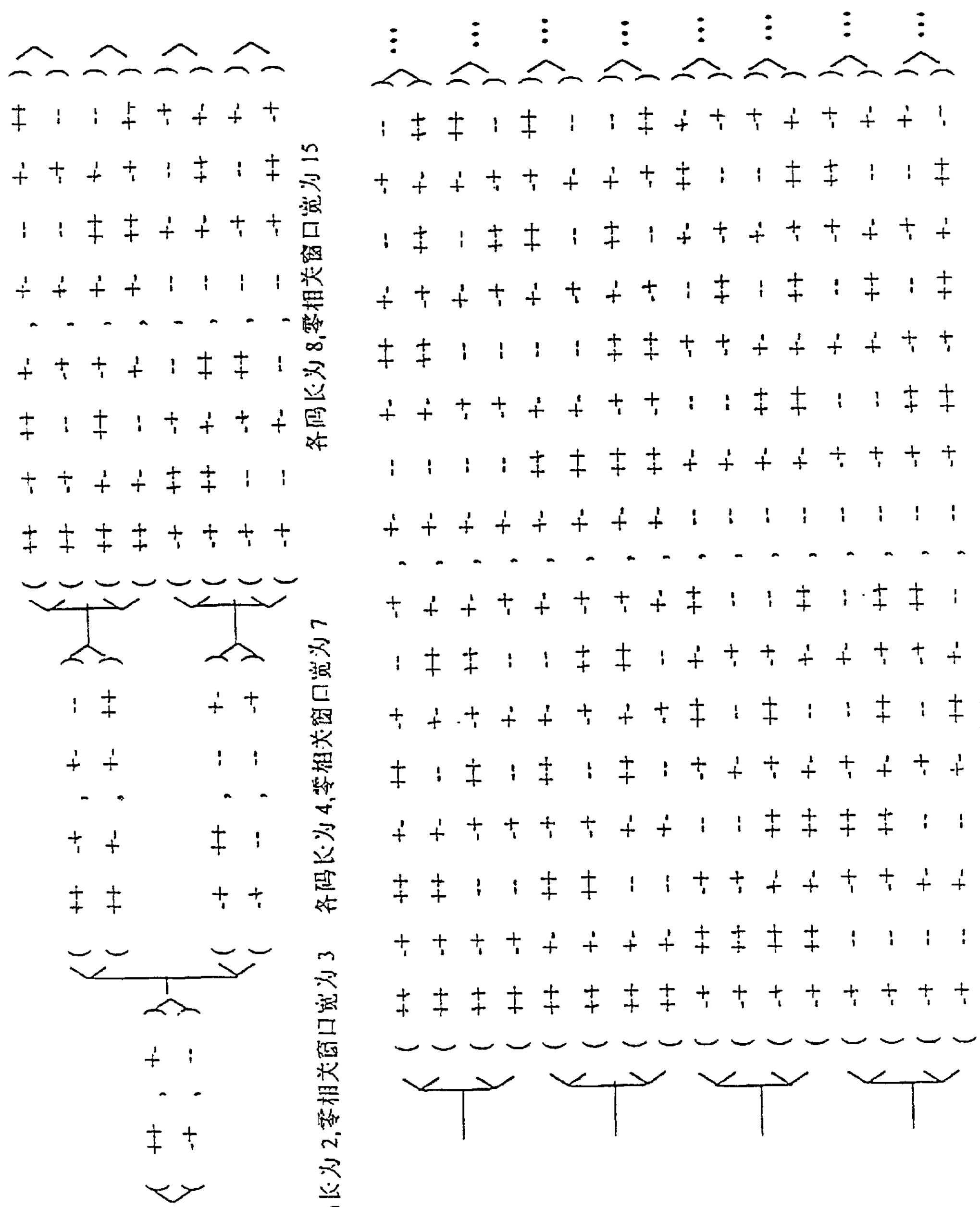
各码长为8,零相关窗口宽为15
各码长为4,零相关窗口宽为7
各码长为2,零相关窗口宽为3
各码长为16,零相关窗口宽为31

图3

Attorney Docket: LKAR01US

DECLARATION AND POWER OF ATTORNEY

As below-named inventor, I hereby declare that: my residence, post office address, and citizenship are as stated below next to my name; that I believe I am the original, first, and sole inventor of the subject matter which is claimed and for which a patent is sought on the invention entitled:

A SPREAD-SPECTRUM MULTIPLE ACCESS CODING METHOD
WITH ZERO CORRELATION WINDOW

the specification for which an international patent application was filed 17 FEBRUARY 2000, having INTERNATIONAL APPLICATION NO. PCT/CN00/00028.

I hereby state that I have reviewed and understand the contents of the above-identified patent application, including the claims.

I acknowledge the duty to disclose information which is material to the examination of this application, in accordance with Title 37, Code of Federal Regulations, § 1.56(a).

I hereby appoint the following attorney to prosecute this application and transact all business in the Patent and Trademark Office connected therewith: David B. Newman, Jr., Registration No. 30,966.

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I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

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